

Unification-Based Grammar Engineering

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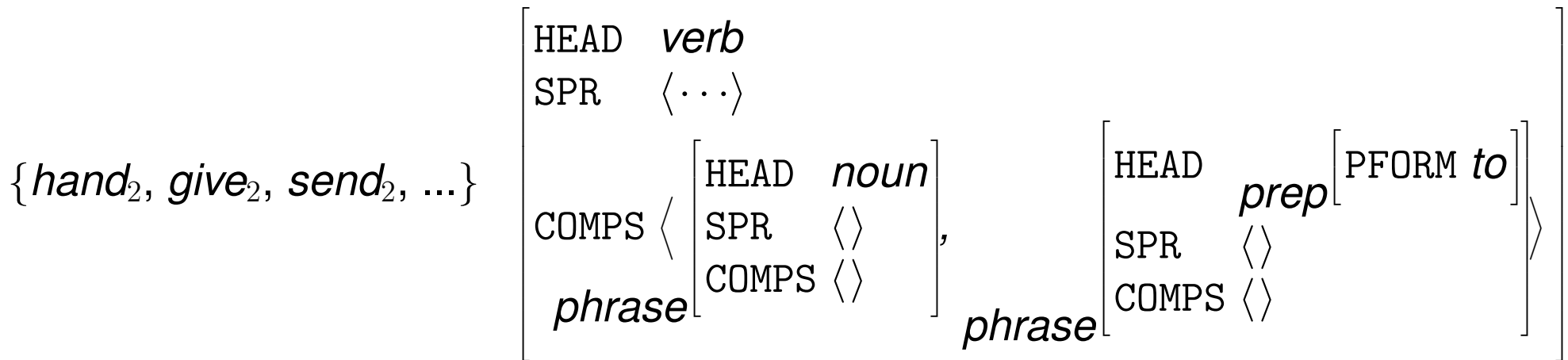
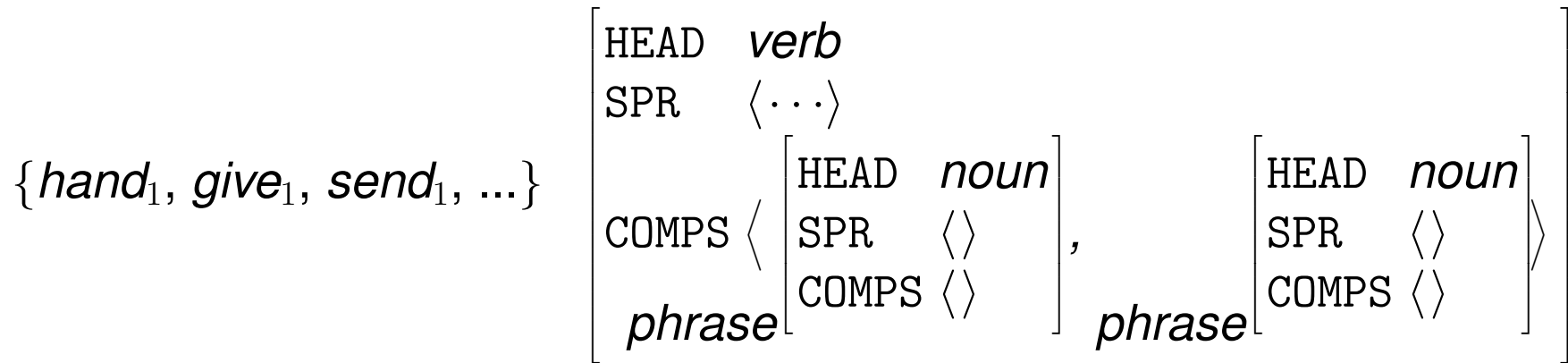
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ESSLLI 2016; August 15–19, 2016

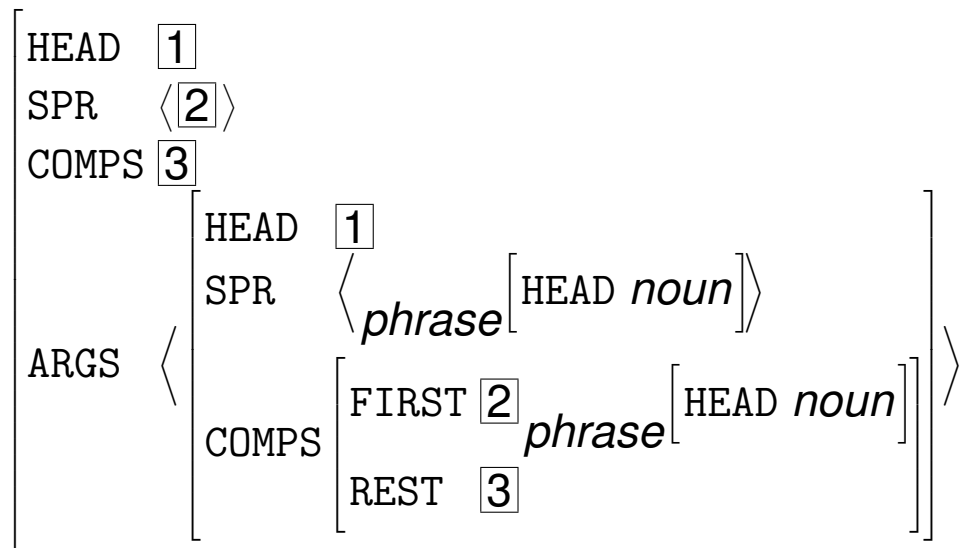
Dative Shift: A Productive Process



Lexical Variation: Lexical Rules

- Dative shift, passivization, et al. are systematic processes in the lexicon;
- use of *monotonic* inheritance is insufficient to relate $give_1$ and $give_2$;
- *lexical rules* are unary grammar rules that operate ‘within the lexicon’;
- take as input a lexical sign (*expression*) and output a derived lexical sign.

Rough Approximation of Passive Lexical Rule



Orthographic Variation: Inflectional Rules

```
%(letter-set (!s abcdefghijklmnopqrstuvwxyz))
```

```
noun-non-3sing_irule :=
```

```
%suffix (!s !ss) (!ss !ssses) (ss sses)
```

```
non-3sing-word &
```

```
[ HEAD [ AGR non-3sing ],
```

```
  ARGS < noun-lxm > ].
```

```
noun-3sing_irule :=
```

```
3sing-word &
```

```
[ ORTH #1,
```

```
  ARGS < noun-lxm & [ ORTH #1 ] > ].
```

dog

|

dogs

bus

|

busses

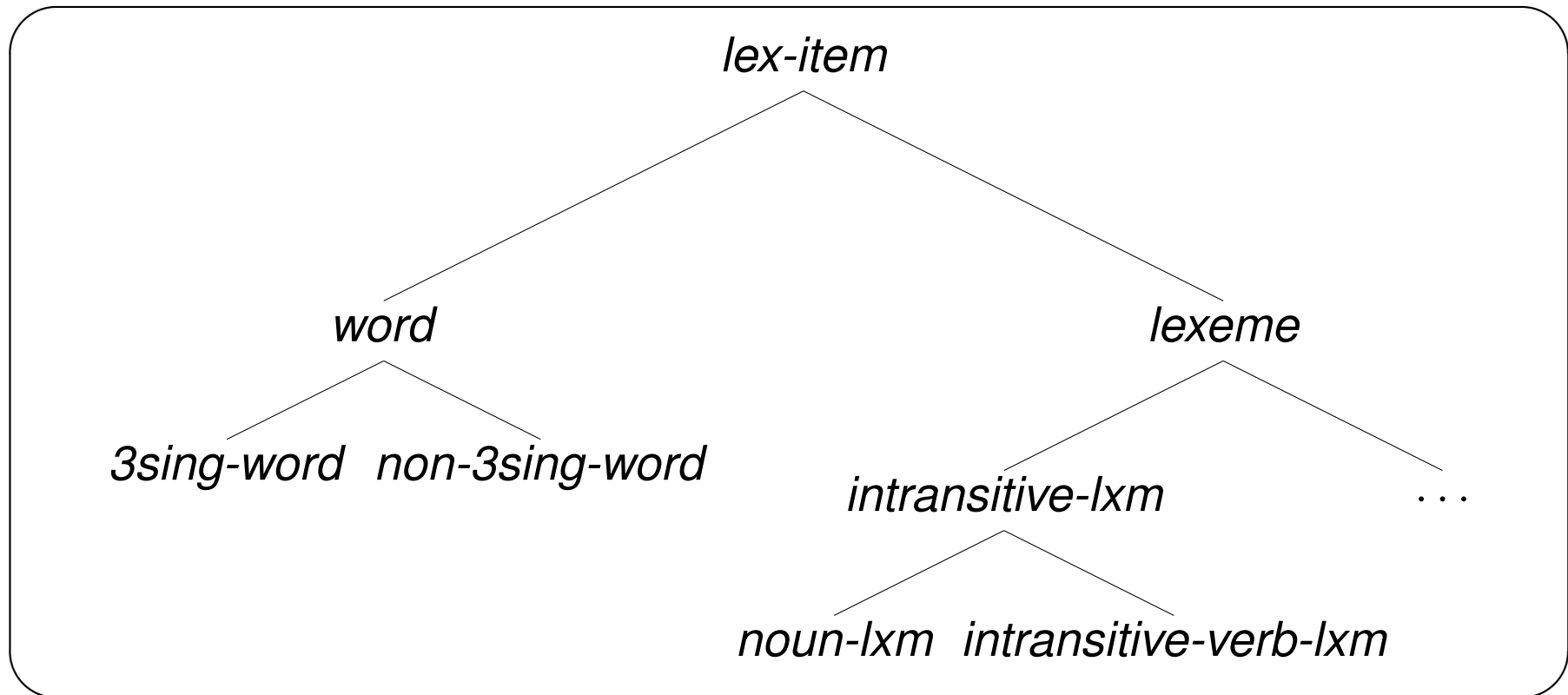
pass

|

passes



The Lexeme vs. Word Distinction



- Lexical entries are *uninflected*; cannot enter syntax by themselves;
- inflectional rules ‘make’ *word* from *lexeme*, possibly with ‘null’ suffix.



Recursion in the Type Hierarchy

- Type hierarchy must be finite *after* type inference; illegal type constraint:

```
*list* := *top* & [ FIRST *top*, REST *list* ].
```

- needs additional provision for empty lists; indirect recursion:

```
*list* := *top*.
```

```
*ne-list* := *list* & [ FIRST *top*, REST *list* ].
```

```
*null* := *list*.
```

- recursive types allow for *parameterized list types* ('list of X'):

```
*s-list* := *list*.
```

```
*s-ne-list* := *ne-list* & *s-list* &  
[ FIRST expression, REST *s-list* ].
```

```
*s-null* := *null* & *s-list*.
```

